Finding Global Minima via Kernel Approximations

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Joint work with Alessandro Rudi and Ulysse Marteau-Ferey Congrès SMAI, la Grande Motte - June 22, 2021

Global optimization

• Zero-th order minimization

 $\min_{x \in \Omega} f(x)$

- $\Omega \subset \mathbb{R}^d$ simple compact subset (e.g., $[-1,1]^d$)
- -f with some bounded derivatives
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- access to function values
- No convexity assumption
- Many applications
 - Hyperparameter optimization in machine learning
 - Industry

Optimal algorithms

- Goal: Find $\hat{x} \in \Omega$ such that $f(\hat{x}) \min_{x \in \Omega} f(x) \leqslant \varepsilon$
 - Lowest number of function calls
 - Worst-case guarantees over all functions f in some convex set ${\mathcal F}$

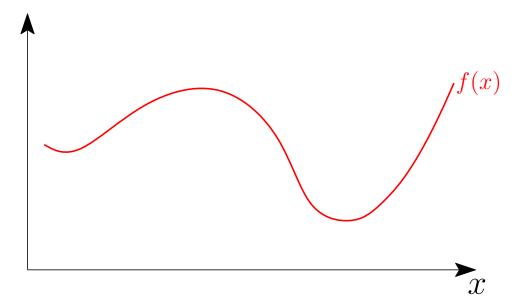
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 - Simplest algorithm: approximate f by \hat{f} and minimize \hat{f}

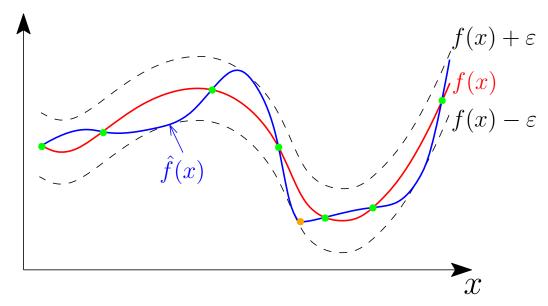


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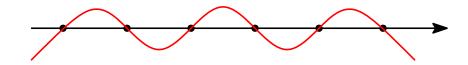
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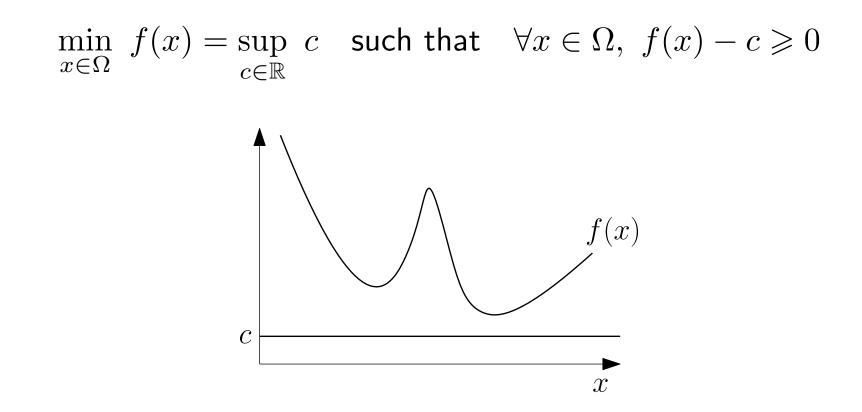


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- Algorithms with polynomial-time complexity in n?
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Reformulations

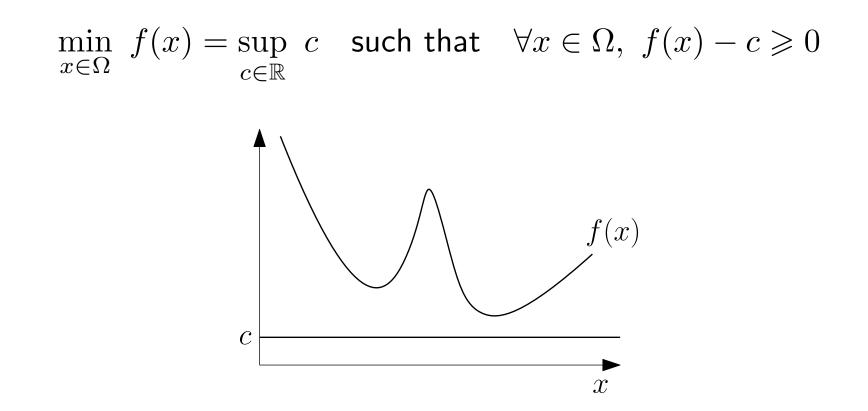
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- Need to represent non-negative functions (such as f(x) c)

Representing non-negative functions

- Assumption: g(x) can be represented as $g(x) = \langle \phi(x), G\phi(x) \rangle$
 - with ${\boldsymbol{G}}$ symmetric operator
 - Assume constant function can be represented as $1=\langle u,\phi(x)\rangle$
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- Positivity through "sums-of-squares"
 - If $G \succcurlyeq 0$, then $\forall x \in \Omega, \ g(x) = \langle \phi(x), G \phi(x) \rangle \ge 0$

- Then,
$$g(x) = \sum_{i \in I} \lambda_i \langle \phi(x), (h_i \otimes h_i) \phi(x) \rangle = \sum_{i \in I} \lambda_i \langle \phi(x), h_i \rangle^2$$

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- Are all non-negative functions sums-of-squares?
 - Polynomials: no if d > 1 (see, e.g., Rudin, 2000)

Global optimization with sums of square polynomials

- Replace $f(x) c \ge 0$ by $f(x) = c + \langle \phi(x), A\phi(x) \rangle$ with $A \succcurlyeq 0$
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• Sum-of-squares optimization (Lasserre, 2001; Parrilo, 2003)

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- Equivalent to original problem if $f(x) - f_*$ is a sum-of-squares - If not, and if localization set $\Omega = \{x, \|x\|^2 \leq R^2\}$ is known,

 $\forall x \in \Omega, \ f(x) \ge 0 \quad \Leftrightarrow \quad \forall x \in \mathbb{R}^d, \ f(x) = q(x) + (R^2 - ||x||^2)p(x)$

with p and q sums-of-squares polynomials (of unknown degree) – Needs "hierarchies"

Representing more general functions with RKHSs

- Reproducing Kernel Hilbert Space (RKHS) :
 - Hilbert space of functions $g \in \mathcal{H}, \ g : \mathbb{R}^d \to \mathbb{R}$
 - Representation as linear form : $g(x) = \langle g, \phi(x) \rangle$
 - Kernel : $k(x, x') = \langle \phi(x), \phi(x') \rangle$ (computable)

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- Example : Sobolev spaces (Berlinet and Thomas-Agnan, 2011)
 - Sobolev spaces $H^s(\Omega)$ with $\Omega \subset \mathbb{R}^d$, s > d/2

$$\langle f,g \rangle = \sum_{|\alpha| \le s} \int_{\Omega} \partial^{\alpha} f(x) \cdot \partial^{\alpha} g(x) dx$$

- Example s = d/2 + 1/2: $k(x, y) = \exp(-||x - y||)$

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\bullet Everything can be expressed using only the kernel function k

- Useful when dealing with function evaluations
- Representer theorem (Kimeldorf and Wahba, 1971): Minimizing $L(g(x_1), \ldots, g(x_n)) + \frac{\lambda}{2} ||g||^2$ can be done by restricting to

$$g(x) = \sum_{i=1}^{n} \alpha_i k(x, x_i)$$

- Then $g(x_j) = \sum_{i=1}^{n} \alpha_i k(x_j, x_i)$ and $\|g\|^2 = \sum_{i,j=1}^{n} \alpha_i \alpha_j k(x_i, x_j)$

Going infinite-dimensional (Rudi, Marteau-Ferey, and Bach, 2020)

 $\sup_{c\in\mathbb{R},\ A\succcurlyeq 0} c \quad \text{such that} \quad \forall x\in\Omega,\ f(x)=c+\langle\phi(x),A\phi(x)\rangle$

- $\phi(x) \in \mathcal{H}$ Hilbert space so that $\langle w, \phi(x) \rangle$ spans a Sobolev space
 - -s > d/2 squared-integrable derivative
 - Reproducing kernel Hilbert space (RKHS)
 - $k(x,y) = \langle \phi(x), \phi(y) \rangle = \exp(-\|x-y\|)$ for s = d/2 + 1/2.
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- Theorem: $\exists A_* \succeq 0$ such that $\forall x \in \Omega$, $f(x) = f_* + \langle \phi(x), A_*\phi(x) \rangle$
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 - \Rightarrow Equivalent to original problem, but infinite-dimensional

• Subsample n points $x_1, \ldots, x_n \in \Omega$ and solve

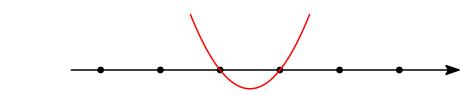
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- Approximation guarantees (Rudi, Marteau-Ferey, and Bach, 2020)
 - With random samples, $n\approx \varepsilon^{-d/(m-d/2-3)}$
 - (up to logarithmic terms)
 - To be compared to optimal rate $n\approx \varepsilon^{-d/(m-d/2)}$
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- Subsampling inequalities as $f(x_i) \ge c$ directly?
 - cannot improve on $n\approx \varepsilon^{-d}$





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- Finite-dimensional algorithm through representer theorem
 - Marteau-Ferey, Bach, and Rudi (2020)
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- Semi-definite programming problem
 - Complexity $O(n^{3.5} \log \frac{1}{\epsilon})$ by interior point method
 - More efficient Newton algorithm in $O(n^3)$

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- 2. Feature computation
 - Compute $K_{ij} = k(x_i, x_j)$ for k Sobolev kernel of smoothness s
 - Compute square root of $K = R^{\top} R \in \mathbb{R}^{n \times n}$
 - Set $\Phi_j = j$ -th column of R, $\forall j \in \{1, \ldots, n\}$
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- 3. Solve $\max_{c \in \mathbb{R}, B \succeq 0} c \lambda \operatorname{tr}(B)$ s.t. $\forall j \in \{1, \dots, n\}, f_j c = \Phi_j^\top B \Phi_j$

– With Lagrange multipliers $\alpha \in \mathbb{R}^n$

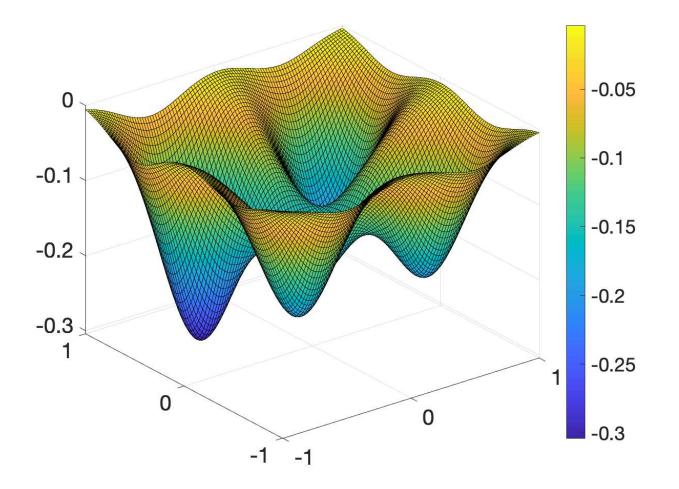
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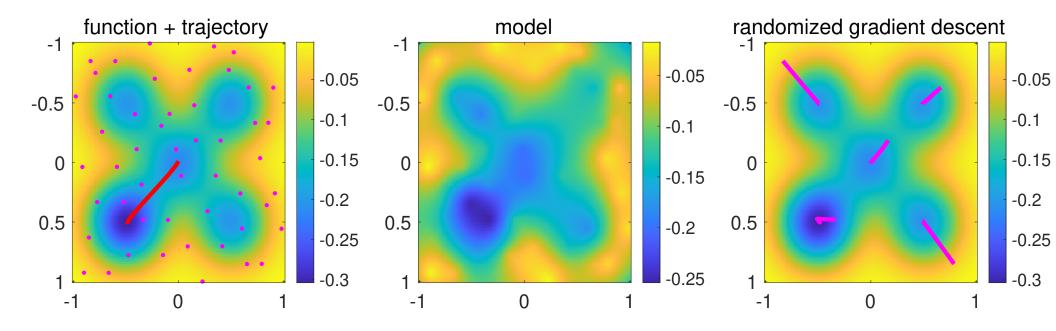
• **Output:** c and $\hat{x} = \sum_{j=1}^{n} \alpha_j x_j$

Illustration

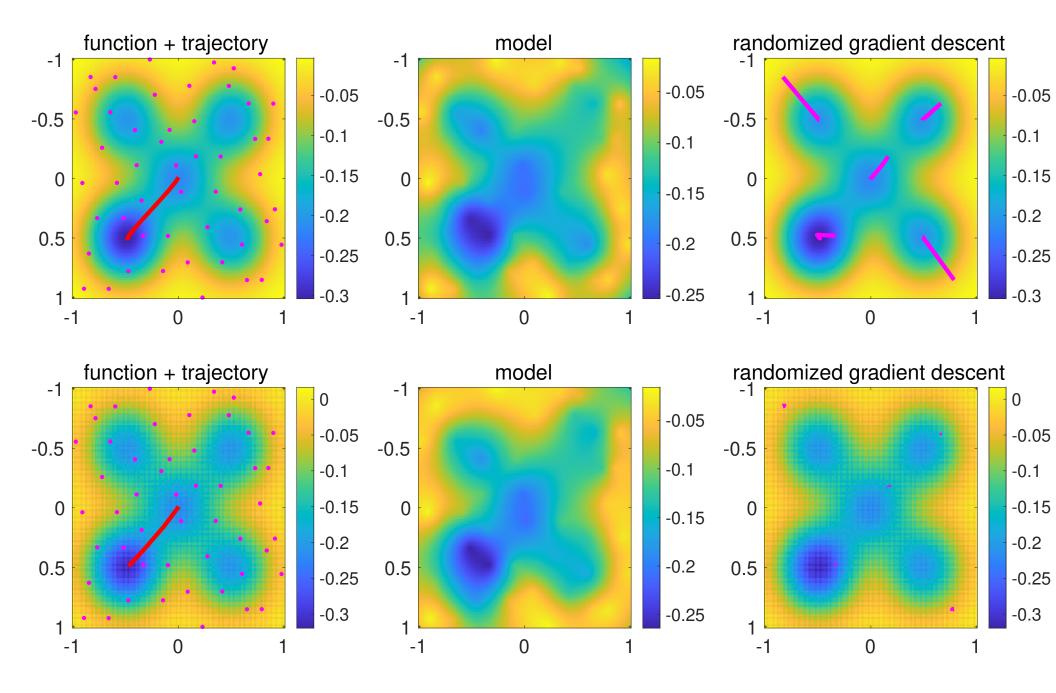
• Minimization of two-dimensional function



Illustration

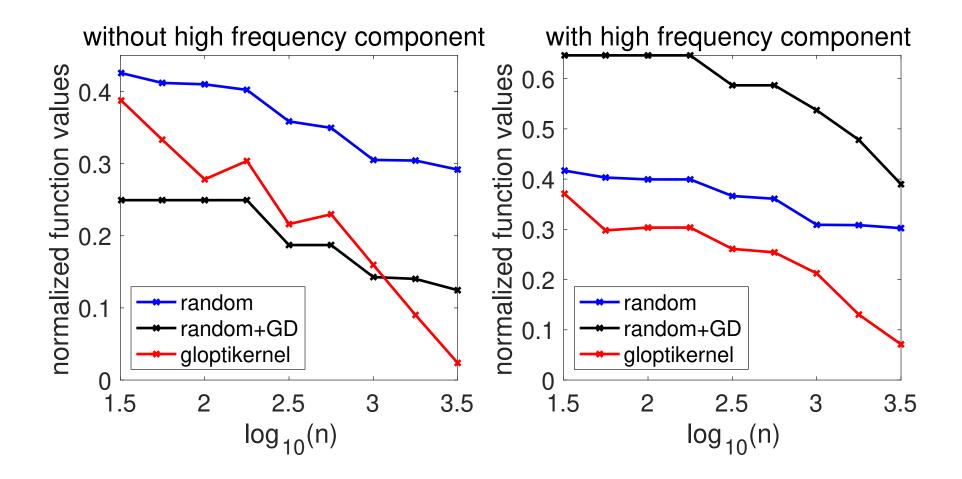


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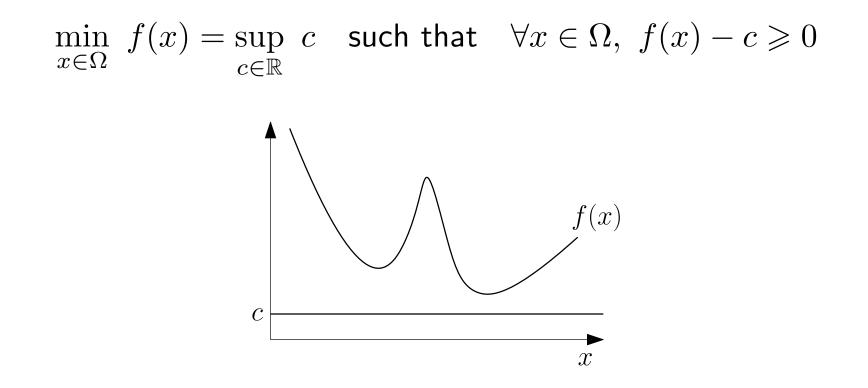
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• Minimization of eight-dimensional function



Duality

• Primal problem



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$$\min_{x\in\Omega} f(x) = \sup_{c\in\mathbb{R}} c \quad \text{such that} \quad \forall x\in\Omega, \ f(x) - c \ge 0$$

• Dual problem on probability measures

$$\inf_{\mu \in \mathbb{R}^{\Omega}} \int_{\Omega} \mu(x) f(x) dx \quad \text{such that} \quad \int_{\Omega} \mu(x) dx = 1, \ \forall x \in \Omega, \ \mu(x) \ge 0$$

Duality with sums-of-squares

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 $\min_{x\in\Omega} f(x) = \sup_{c\in\mathbb{R},\ A\succcurlyeq 0} c \quad \text{such that} \quad \forall x\in\Omega,\ f(x) - c = \langle \phi(x), A\phi(x) \rangle$

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- Extension of results on polynomials (Lasserre, 2020)

Extension - I

• Generic constrained optimization problem

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- Application to optimal transport (Vacher, Muzellec, Rudi, Bach, and Vialard, 2021)

Smooth optimal transport (Vacher et al., 2021)

- Primal formulation: $\inf_{\gamma \in \Gamma(\mu,\nu)} \int_{\mathfrak{X} \times \mathfrak{Y}} c(x,y) d\gamma(x,y)$
 - $\Gamma(\mu,\nu)$ set of probability distributions with marginals μ and ν
- Dual formulation: $\sup_{u,v\in C(\mathbb{R}^d)} \int_{\mathcal{X}} u(x)d\mu(x) + \int_{\mathcal{Y}} v(y)d\mu(y)$ such that $\forall (x,y) \in \mathcal{X} \times \mathcal{Y}, \ c(x,y) \ge u(x) + v(y)$

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- Kernel sums of squares: replace inequality constraint by: $\forall (x,y) \in \mathfrak{X} \times \mathfrak{Y}, \ c(x,y) = u(x) + v(y) + \langle \phi(x,y), A\phi(x,y) \rangle$

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Conclusion

• Global optimization through kernel approximations

- Joint optimization and approximation
- infinite-dimensional sums-of-squares representation
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• Further extensions

- Efficient algorithms below ${\cal O}(n^3)$ complexity
- Adaptive choice of sampling points
- Certificates of optimality
- Other infinite-dimensional convex optimization problems

Conclusion

• Global optimization through kernel approximations

- Joint optimization and approximation
- infinite-dimensional sums-of-squares representation
- Controlled subsampling with guarantees

• Further extensions

- Efficient algorithms below ${\cal O}(n^3)$ complexity
- Adaptive choice of sampling points
- Certificates of optimality
- Other infinite-dimensional convex optimization problems
- See arxiv.org/abs/2012.11978 and francisbach.com/
- See talk by Ulysse Marteau-Ferey (Wednesday at 11am)

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